

# Admin

- ◇ Assign 2 due Wed
- ◇ Today's topics
  - Functional recursion
- ◇ Reading
  - Reader ch. 4-5-6 (today-W-F)

Lecture #8

# Solving problems recursively

- ◇ A recursive function calls itself — wacky!
- ◇ Idea: solve problem using coworkers (clones) who work and act like you
  - Delegate similar, smaller problem to clone
  - Combine result from clone(s) to solve total problem
  - Work toward trivial version that is directly solvable
- ◇ For problems that exhibit "self-similarity"
  - Structure repeats within at different levels of scale
  - Solving larger problem means solving smaller problem(s) within
- ◇ Feels mysterious at first
  - "Leap of faith" required
  - With practice, master the art of recursive decomposition
  - Eventually grok the underlying patterns

# Functional recursion

- ◇ Function that returns answer/result
  - Outer problem result uses result from smaller, same problem(s)
- ◇ Base case
  - Simplest version of problem
  - Can be directly solved
- ◇ Recursive case
  - Make call(s) to self to get results for smaller, simpler version(s)
  - Recursive calls must advance toward base case
  - Results of recursive calls combined to solve larger version

# Power example

- ◇ C++ has no exponentiation op
- ◇ Iterative formulation for Raise function
  - $\text{base}^{\text{exp}} = \text{base} * \text{base} * \dots * \text{base}$  (exp times)

```
int Raise(int base, int exp)
{
    int result = 1;
    for (int i = 0; i < exp; i++)
        result *= base;
    return result;
}
```

## Recursive version

### ◇ Now consider recursive formulation

- $\text{base}^{\text{exp}} = \text{base} * \text{base}^{\text{exp}-1}$

```
int Raise(int base, int exp)
{
    if (exp == 0) return 1; } Base case
    else
        return base * Raise(base, exp-1); } Recursive case
}
```

## More efficient recursion

$\text{base}^{\text{exp}} = \text{base}^{\text{exp}/2} * \text{base}^{\text{exp}/2}$  (\* base if exp is odd)

```
int Raise(int base, int exp)
{
    if (exp == 0)
        return 1;
    else {
        int half = Raise(base, exp/2);
        if (exp % 2 == 0)
            return half * half;
        else
            return base * half * half;
    }
}
```

## Avoid "arm's length" recursion

### ◇ Aim for simple, clean base case

- No need to anticipate other earlier stopping points
- Avoid looking ahead before recursive calls, just let simple base case handle

```
int Raise(int base, int exp)
{
    if (exp == 0) return 1;
    else if (exp == 1) return base;
    else if (exp == 2) return base * base;
    else if (exp == 3) return base * base * base;
    else return base * Raise(base, exp - 1);
}
```

## Recursion and efficiency

- ◇ Recursion provides no guarantee of (in)efficiency
  - Recursion can require same resources as alternative approach
    - Or recursion may be much more or much less efficient
  - For problems with simple iterative solution, iteration is likely the best
- ◇ Why recursion then?
  - Can express with clear, direct, elegant code
  - Can intuitively model a task that is recursive in nature
  - Solution may require recursion — iteration won't do!

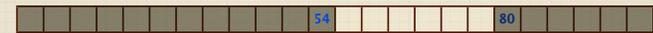
# Palindromes

- ◇ A palindrome string reads same when reversed
  - e.g. "was it a car or a cat i saw", "go hang a salami im a lasagna hog"
- ◇ Recursive insight
  - First and last letter match and interior is palindrome
- ◇ Base case?

```
bool IsPalindrome(string s)
{
    if (s.length() <= 1) return true;
    return s[0] == s[s.length()-1] &&
        IsPalindrome(s.substr(1, s.length()-2));
}
```

# Binary search

- ◇ Searching for key within vector
  - *Linear search* starts at beginning and searches to end
  - *Binary search* uses divide-and-conquer (requires **sorted** vector)
    - Much faster method!
- ◇ Recursive insight:
  - Consider middle elem of vector, if key, you're done
  - Otherwise decide which half to recursively search



- ◇ Base case?

# Binary search code

```
const int NotFound = -1;

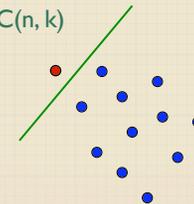
int BSearch(Vector<string> &v,
            int start, int stop, string key)
{
    if (start > stop) return NotFound;

    int mid = (start + stop)/2;
    if (key == v[mid])
        return mid;
    else if (key < v[mid])
        return BSearch(v, start, mid-1, key);
    else
        return BSearch(v, mid+1, stop, key);
}
```

- ◇ Classic "divide and conquer" algorithm
  - Operates very efficiently! Double size of vector, how much longer to search?

# Choosing a subset

- ◇ Reader ch 4, exercise 8
  - Given N things, how many different ways can you choose K of them?
    - e.g. given a dorm of 60 people, how many different groups of 4 people can go together to Flicks?
  - N-choose-K, written as  $C(n, k)$



Number of subsets that include ● =  $C(n-1, k-1)$   
+ Number of subsets that don't include ● =  $C(n-1, k)$

# Choose code

## ◇ Simplest base case

- when **no** choices remain at all

```
int C(int n, int k)
{
    if (k == 0 || k == n)
        return 1;
    else
        return C(n-1, k) + C(n-1, k-1);
}
```